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Logic Exercise Sheet 6

Exercise 6.1 Validity and satisfiability

Which of the following formulas is valid, which is satisfiable?

For each satisfiable formula give a model. For each formula that is satisfiable but not valid, give a counter-model.

- a) $(\forall x. \exists y. P(x, y)) \Rightarrow (\exists y. \forall x. P(x, y))$
- b) $(\exists x. \forall y. P(x, y)) \Rightarrow (\forall y. \exists x. P(x, y))$
- c) $(\exists x. \forall y. P(x, y)) \Rightarrow (\forall x. \exists y. P(x, y))$

Exercise 6.2 More validity, satisfiability and models

Prove or disprove, using the definitions:

- a) For each sentence p, either p is valid or $\neg p$ is valid.
- b) A sentence p is valid iff $\neg p$ is unsatisfiable.

Exercise 6.3 Trivial model

Prove that a formula containing only $\land, \lor, \forall, \exists, \Rightarrow$ and atomic formulas is always satisfiable.

Exercise 6.4 Extension with \exists !

- a) Extend the definition of first-order formulas with an additional quantifier $\exists !,$ such that $\exists !x.P(x)$ means that there exists a unique x such that P(x) holds. Modify the definition of the semantics.
- b) Extend the OCaml implementations of the type formula and the function holds accordingly.
- c) Is the following formula valid?

 $(\exists !x.\exists !y.P(x,y)) \Rightarrow (\exists !y.\exists !x.P(x,y))$